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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

	Application No.	Applicant(s)
Office Action Summary	10/734,028	DILLAWAY ET AL.
	Examiner	Art Unit	
	CARLTON V. JOHNSON	2436	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE _3_ MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.
- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133).
  Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

1)☒ Responsive to communication(s) filed on _21 October 2009_.

2a)☐ This action is **FINAL**. 2b)☒ This action is non-final.

3)☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

4)☒ Claim(s) _2-4,7-12,14,16,18-21,23,25-33 and 36-38_ is/are pending in the application.

    4a) Of the above claim(s) _____ is/are withdrawn from consideration.

5)☐ Claim(s) _____ is/are allowed.

6)☒ Claim(s) _2-4,7-12,14,16,18-21,23,25-33 and 36-38_ is/are rejected.

7)☐ Claim(s) _____ is/are objected to.

8)☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

9)☐ The specification is objected to by the Examiner.

10)☐ The drawing(s) filed on _____ is/are: a)☐ accepted or b)☐ objected to by the Examiner.

    Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).

    Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).

11)☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

12)☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).

    a)☐ All b)☐ Some * c)☐ None of:

      1.☐ Certified copies of the priority documents have been received.

      2.☐ Certified copies of the priority documents have been received in Application No. _____.

      3.☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

    * See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

1)☒ Notice of References Cited (PTO-892)

2)☐ Notice of Draftsperson's Patent Drawing Review (PTO-948)

3)☐ Information Disclosure Statement(s) (PTO/SB/08)
    Paper No(s)/Mail Date _____.

4)☐ Interview Summary (PTO-413)
    Paper No(s)/Mail Date. _____ .

5)☐ Notice of Informal Patent Application

6)☐ Other: _____.

DETAILED ACTION

Continued Examination Under 37 CFR 1.114

1.   A request for continued examination under 37 CFR 1.114, including the fee set forth in 37 CFR 1.17(e), was filed in this application after final rejection. Since this application is eligible for continued examination under 37 CFR 1.114, and the fee set forth in 37 CFR 1.17(e) has been timely paid, the finality of the previous Office action has been withdrawn pursuant to 37 CFR 1.114.

Applicant's submission filed on 10-21-2009 has been entered.

2.   Claims **2 - 4, 7 - 12, 14, 16, 18 - 21, 23, 25 - 33, 36 - 38** are pending.   Claims **11, 19, 25 - 27, 29 - 31, 33** have been amended.   Claims **37, 38** are new. Claims **1, 5, 6, 13, 15, 17, 22, 24, 34, 35** have been cancelled.   Claims **19, 30, 31** are independent. This application was filed on 12-11-2003.

Response to Arguments

3.   Applicant's arguments have been fully considered and they were persuasive and new grounds of rejection have been entered.

3.2   Applicant argues that the referenced prior art does not disclose, *two step process for attestation.*

     Yan prior art discloses that there are two message flows used for the initiation and completion of attestation between two entities. The two message-flows (paragraph

[0064], lines 1-4; paragraph [0064], lines 8-10) initiate attestation and establish the

requirements for the attestation information transferred between the entities. (see Yan

paragraph [0064], lines 1-17: initial attestation transmitted via message flow between

trustor and trustee to establish trust relationship; Yan discloses a two step process for

attestation messages (message flows 302, 304 and message flows 306, 308))

The first message flow is 302 and 304 in Figure 3 of Yan discloses the initial attestation

between the trustor and trustee and an attestation initiation protocol sequence. The

Examiner feels these message flows are analogous to the can-attest and attestation-

wanted request/response sequence. These message flows initiate an attestation

process. Message flows 306 and 308 are used to establish the trust condition between

the trustor and trustee.

Specification in paragraph [0053] discloses that the can-attest message (request)

can have any format and contain any relevant information such as identification of first

entity. Message flows 302 and 304 should satisfy this requirement. Specification in

paragraph [0055] discloses that the attestation-wanted message (response) contains

the requirements for the attestation. Message flows 306 and 308 should satisfy this

requirement.

The claimed invention indicates a first entity and a second entity. The can-attest

message can have any format and contain any information. It appears that a first entity

can send the initial request in a situation (Specification paragraph [0052]). The Pinkas

prior art discloses a situation where a trustee sends an initial query message for

attestation and receives the attestation requirements in response. This message flow is

not novel.

Claim Rejections - 35 USC § 112

4.   The following is a quotation of the first paragraph of 35 U.S.C. 112:

> The specification shall contain a written description of the invention, and of the manner and process of making and using it, in such full, clear, concise, and exact terms as to enable any person skilled in the art to which it pertains, or with which it is most nearly connected, to make and use the same and shall set forth the best mode contemplated by the inventor of carrying out his invention.

5.   Claims **37, 38** are rejected under 35 U.S.C. 112, first paragraph, as failing to

comply with the written description requirement. The claim(s) contains subject matter

which was not described in the specification in such a way as to reasonably convey to

one skilled in the relevant art that the inventor(s), at the time the application was filed,

had possession of the claimed invention.

There does not appear to be any disclosure in the specification for the Claim **37**

limitation: *"the can-attest message is transmitted from the first computer entity to the*

server in an unencrypted format", in the specification or original claims. This will be

considered new matter.   The term, *"unencrypt"*, only appears in the specification in

paragraph [0043] in reference to the transmission of a symmetric key.

There does not appear to be any disclosure in the for the Claim **38** limitation:

"each of the can-attest message and the attestation-wanted message is configured to

preclude a) containing or b) using a cryptographic key".   The term "preclude" does not

appear in the specification or original claim.   The specification does not appear to

indicate that the can-attest message and attestation-wanted messages cannot contain

or use a cryptographic key.

Appropriate correction required.

Specification paragraph [0043]:
> [0043] For example, and as seen in Fig. 2, the trust message 24 may include therein a symmetric key (K) that the first and second entities 10, 12 shall each employ to encrypt and decrypt messages therebetween. **Of course, inasmuch as such a symmetric key (K) should not be transmitted to the first entity 10 in an unencrypted format,** the second entity 12 should encrypt (K) according to a key decryptable by the first entity 10. ....

Claim Rejections - 35 USC § 103

4.   The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all

obviousness rejections set forth in this Office action:

> a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negatived by the manner in which the invention was made.

5.   Claims **2 - 4, 7 - 12, 14, 18 - 21, 23, 25 - 27, 29, 31 - 37** are rejected under 35

U.S.C. 103(a) as being unpatentable over **Yan et al.** (US PGPUB No. **20050033987**) in

view of **Qui** (US PGPUB No. **20040148505**) and further in view of **Pinkas et al.** (US

Patent No. **5,214,700**).

Regarding Claim 31, Yan discloses a method of establishing trust between two

computer entities, the method comprising:

a) transmitting an attestation message from a first computer entity to a second

computer entity, the attestation message including a code identifier (code ID) that

is calculated by using a security ID corresponding to a behavior parameter that is

associated with a computing operation having security implications; (see Yan

paragraph [0064], lines 1-17: initial attestation transmitted via message flow

between trustor and trustee to establish trust relationship; Yan discloses a two

step process for attestation messages (message flows 302, 304 and message

flows 306, 308))

Furthermore, Yan discloses:

b) ensuring that the security ID corresponding to the behavior parameter has not

    been tampered with <u>by verifying the validity of the code ID in the second</u>

    <u>computer entity, the verifying comprising determining that the first computer</u>

    <u>entity is not included in a do-not-trust list</u>; and transmitting a trust message from

    the second computer entity to the first computer entity upon successfully verifying

    the validity of the code ID, the trust message including a first secret that is shared

    between the first and the second computer entities for communicating securely.

    (see Yan paragraph [0060], lines 6-9: verify signature, attestation information;

    paragraph [0062], lines 1-4; paragraph [0064], lines 1-4: session key exchanged

    between entities for future messaging, communications; paragraph [0060], lines

    6-9: check certificate in certificate chain, not on revoked list (do-not-trust list))

Yan does not specifically disclose a period time whereby a secret is valid.

However, Qui discloses a first period of time, wherein the first period of time is

determined by the second computer entity. (see Qui paragraph [0040], lines 1-7;

paragraph [0021], lines 8-11: expiration timer for certificate information)

It would have been obvious to one of ordinary skill in the art to modify Yan for a period time whereby a secret is valid as taught by Qui.   One of ordinary skill in the art would have been motivated to employ the teachings of Qui in order to enable the capability for the generation, transmission, and updating of certificate information when the number of devices is large. (see Qui paragraph [0007], lines 7-12)

Yan does not specifically disclose a first entity sends an initial attestation. Pinkas discloses or can-attest message or an initial message transmitted from a first entity. (see Pinkas col 4, line 54 - col. 5, line 2: requestor (trustee, first entity) addresses an attestation request or query to the authority; server transmits attestation to requestor subject)

It would have been obvious to one of ordinary skill in the art to modify Yan-Qui whereby a first entity sends an initial attestation message as taught by Pinkas.   One of ordinary skill in the art would have been motivated to employ the teachings of Pinkas to provide flexible and simplicity in the establishment of attestation while provide security with the usage of a check datum. (see Pinkas col. 4, lines 44-53)

Regarding Claims 2, Yan discloses the method of claim 31 wherein the first computer entity encrypts the code ID of the attestation message according to a key available to the second computer entity, the method further comprising the second computer entity decrypting such encrypted matter. (see Yan paragraph [0059], lines 6-12: certificate (public/private) key available to second entity, utilized to encrypt (signature) attestation

information)

Regarding Claims 3, Yan discloses the method of claim 31 wherein the second

computer entity consumes the attestation message by application of same to a verifying

function that automatically verifies the attestation message based on a format thereof

and that extracts relevant information from such verified attestation message for use by

the second computer entity. (see Yan paragraph [0060], lines 6-9: verify attestation

message with extracted information based on formatted information (certificate

information, encrypted hash))

Regarding Claims 4, Yan discloses the method of claim 31 wherein the first computer

entity is a part of a computing device, and the second computer entity decides based on

the code ID in the attestation message whether the first computer entity can be trusted,

and also decides based on a certificate chain of the message whether the computing

device can be trusted, the certificate chain leading back to a trusted root authority. (see

Yan paragraph [0060], lines 9-12: entity can be trusted; paragraph [0060], lines 1-6:

certificate chain utilized to establish trust)

Regarding Claims 7, Yan discloses the method of claim 4 wherein the second

computer entity determines that the code ID is a known code ID and that the first

computer entity can be trusted based on such code ID. (see Yan paragraph [0060],

lines 6-9: check first entity on application trust list based on integrity metric (code ID) of

first entity)

Regarding Claims 8, Yan discloses the method of claim 4 wherein the second

computer entity determines from the certificate chain whether the computing device of

the first computer entity should be trusted to instantiate and operate the first computer

entity in a trusted manner and should be trusted to calculate the code ID properly. (see

Yan paragraph [0060], lines 1-9: certificate information utilized to determine trust status)

Regarding Claims 9, Yan discloses the method of claim 8 wherein the second

computer entity determines that each certificate in the certificate chain is not on a do-

not-trust list. (see Yan paragraph [0060], lines 6-9: check certificate in certificate chain,

not on revoked list (do-not-trust list))

Regarding Claims 10, Yan discloses the method of claim 31 wherein the trust

message includes a symmetric key (K) that the first and second computer entities shall

each employ to encrypt and decrypt messages therebetween. (see Yan paragraph

[0062], lines 1-4; paragraph [0064], lines 1-4: session key exchanged between entities

for future messaging)

Regarding Claims 11, Yan discloses the method of claim 10 wherein the symmetric

key (K) is encrypted according to a public key of the first entity (PU-1) to result in (PU-

1(K)), the second entity obtaining (PU-1) from the certificate chain of the attestation

message, and wherein the first computer entity obtains the symmetric key (K) from the

received trust message by applying a private key (PR-1) corresponding to (PU-1) to

(PU-1(K)) to result in (K). (see Yan paragraph [0059], lines 6-12: public/private

certificate key, encrypt (signature) attestation information)

Regarding Claims 12, Yan discloses the method of claim 31 wherein the trust

message further includes an identification of a cryptographic algorithm to be employed

in connection with the first secret. (see Yan paragraph [0062], lines 1-4; paragraph

[0064], lines 1-4: session key for secure communications (interaction); paragraph

[0065], lines 9-15: updated attestation information (cryptographic algorithm), protocol for

exchange negotiated)

Regarding Claims 14, Yan discloses the method of claim 31 wherein the trust

message further includes relevant trust data encrypted according to a key available to

the first computer entity, and wherein the first computer entity decrypts the encrypted

trust data by applying the key thereto. (see Yan paragraph [0059], lines 6-12:

public/private key certificate (available to first entity) used to encrypt (signature)

attestation information)

Regarding Claims 18, Yan discloses the method of claim 31 wherein prior to the first

computer entity transmitting the attestation message, the first computer entity sends a

can-attest message to the second computer entity, the can-attest message stating that

the first computer entity can send an attestation message but that the first computer

entity would like to know from the second computer entity whether such an attestation

message is required by such second computer entity and if so any requirements that

such second computer entity has with regard to such attestation message, the method

further comprising the second computer entity sending an attestation-wanted message

to the first computer entity in response to the can-attest message, the attestation-

wanted message stating that the second computer entity does in fact require an

attestation message from the first computer entity and that the attestation message as

sent by the first computer entity must adhere to certain requirements as defined in such

attestation-wanted message, whereby the first computer entity thereafter sends the

attestation message in accordance with the requirements stated in the attestation-

wanted message.   (see Yan paragraph [0064], lines 1-17: initial attestation transmitted

via message flow between trustor and trustee to establish trust relationship; Yan

discloses a two step process for attestation messages (message flows 302, 304 and

message flows 306, 308)); paragraph [0065], lines 9-15: update attestation (wanted-

message) information, negotiate protocol for exchange of attestation information)

Yan does not specifically disclose a first entity sends an initial attestation.

However, Pinkas discloses or can-attest message or an initial message transmitted

from a first entity. (see Pinkas col 4, line 54 - col. 5, line 2: requestor (trustee, first

entity) addresses an attestation request or query to the authority; server transmits

attestation to requestor subject)

        It would have been obvious to one of ordinary skill in the art to modify Yan-Qui

whereby a first entity sends an initial attestation message as taught by Pinkas.   One of

ordinary skill in the art would have been motivated to employ the teachings of Pinkas to

provide flexible and simplicity in the establishment of attestation while provide security

with the usage of a check datum. (see Pinkas col. 4, lines 44-53)

Regarding Claims 19, Yan discloses the method of claim 30 further comprising:

a) the first computer entity constructing, in accordance with the requirements stated

   in the attestation-wanted message, the attestation message to be delivered to

   the <u>server</u>, the attestation message including a code identifier (code ID)

   representative of the first computer entity and data relevant to the purpose of the

   trust-based relationship; (see Yan paragraph [0064], lines 1-17: initial attestation

   transmitted via message flow between trustor and trustee to establish trust

   relationship; Yan discloses a two step process for attestation messages

   (message flows 302, 304 and message flows 306, 308); paragraph [0060], lines

   6-9: check entity integrity metric (code ID), identify on application trust list)

Furthermore, Yan discloses the following:

b) the first computer entity appending a digital signature to the attestation message

   and a certificate chain leading back to a trusted root authority, the signature

   being based on the code ID and data thereof and being verifiable based on a

   security key included in the certificate chain, the certificate chain including at

   least one certificate therein proffering trustworthiness of the first computer entity;

   (see Yan paragraph [0060], lines 1-6: certificate chain utilized for attestation

   information, (exchange, verification))

c) the first entity sending the attestation message to the <u>server</u> and the second

entity receiving same, whereby the <u>server</u> verifies the signature of the received

attestation message based on the included security key (see Yan paragraph

[0060], lines 6-9: verify signature, attestation information), whereby alteration of

the code ID or data of the attestation message should cause the signature to fail

to verify, the <u>server</u> based on such a failure dishonoring such attestation

message, the second entity decides whether to in fact enter into the trust-based

relationship with the first entity based on the code ID and the data in the

attestation message, the <u>server</u> upon deciding to in fact enter into the trust-based

relationship with the first entity constructs a trust message to be delivered to the

first entity, the trust message establishing the trust-based relationship and

including therein a secret to be shared between the first <u>entity and the server</u>,

where such shared secret allows such first <u>entity and the server</u> to communicate

in a secure manner, and the <u>entity</u> sends the trust message to the first entity and

the first entity receiving same; (see Yan paragraph [0062], lines 1-4; paragraph

[0064], lines 1-4: session key exchanged between entities for future messaging,

communications; paragraph [0060], lines 6-9: check entity integrity metric (code

ID), identify on application trust list)

d) the first entity obtaining the shared secret in the trust message and employing the

shared secret to exchange information with the <u>server</u> according to the

established trust-based relationship with such <u>server</u>. (see Yan paragraph

[0062], lines 1-4; paragraph [0064], lines 1-4: session key exchanged between

entities for future messaging, communications)

Regarding Claims 20, Yan discloses the method of claim 19 wherein the code identifier (code ID) is calculated from the digest of the first computer entity, whereby alteration of the first computer entity causes the code ID to change. (see Yan paragraph [0065], lines 9-15: updated integrity metric (code ID) modified, new attestation protocol required)

Regarding Claims 21, Yan discloses the method of claim 20 wherein the code identifier (code ID) is calculated from the digest of the first computer entity and from security information relating thereto, whereby alteration of the first computer entity or the security information causes the code ID to change. (see Yan paragraph [0065], lines 9-15: alteration of security information causes integrity metric (code ID) to be modified)

Regarding Claims 23, Yan discloses the method of claim 19 further comprising a code ID calculator of the first computer entity that is used for calculating the code ID, the code ID calculator operating in a trusted manner in a computing device. (see Yan paragraph [0020], lines 7-8: generate integrity metric (code ID) on trusted device)

Regarding Claims 25, Yan discloses the method of claim 19 wherein the first computer entity creates the attestation message by application of the code ID and data thereof to a quoting function that automatically produces the attestation message in an

appropriate format that is accessible to the <u>server</u>. (see Yan paragraph [0065], lines 9-15: attestation information generated in an accessible format (negotiated) with second entity))

Regarding Claims 26, Yan discloses the method of claim 19 wherein the <u>server</u> constructs a trust message including therein a shared secret comprising a symmetric key (K) that the first <u>entity and the server</u> employ to encrypt and decrypt messages therebetween, the symmetric key (K) being encrypted according to a public key (PU-1) to result in (PU-1(K)), the <u>server</u> obtaining (PU-1) from the certificate chain of the attestation message, the method comprising the first entity obtaining the symmetric key (K) from the received trust message by applying a private key (PR-1) corresponding to (PU-1) to (PU-1(K)) to result in (K). (see Yan paragraph [0062], lines 1-4; paragraph [0064], lines 1-4: session key utilized for messaging between first and second entities; paragraph [0059], lines 4-9: public/private key for encryption/decryption (signature attachment))

Regarding Claims 27, Yan discloses the method of claim 19 wherein the <u>server</u> constructs a trust message further including relevant trust data encrypted according to a key available to the first computer entity, the method comprising the first computer entity decrypting the encrypted trust data by applying the key thereto. (see Yan paragraph [0059], lines 6-12: attestation information encrypted (signature) based on public/private keys (known to first entity))

Regarding Claims 29, Yan discloses the method of claim 19 whereby the trust

message is a first trust message and the shared secret is a first shared secret, and

whereby the <u>server</u> constructs a second trust message to be delivered to the first

computer entity, the second trust message including therein a second secret to be

shared between the first <u>entity and the server</u>, where such second shared secret allows

such first <u>entity and the server</u> to communicate in a secure manner (see Yan paragraph

[0062], lines 1-4; paragraph [0064] lines 1-4: session key exchanged between entities

for future messaging), and the second computer entity sends the second trust message

to the first computer entity and the first computer entity receives same, the method

further comprising the first computer entity obtaining the second shared secret in the

trust message and employing the second shared secret to exchange information with

the <u>server</u>, whereby the first shared secret is no longer valid. (see Yan paragraph

[0065], lines 9-15: update attestation information, previous attestation information

invalid)

Regarding Claim 32, Yan discloses the method of claim 31, wherein the security ID is

stored in a location in the first computer entity, and wherein the first computer entity is

constrained to executing a particular behavior only via accessing the stored location.

(see Yan paragraph [0040], lines 10-13: provides a facility whereby a platform may

store secrets accessible only when platform is in a defined configuration)

Regarding Claim 33, Yan discloses the method of claim 31, wherein the behavior

parameter comprises <u>at least one of a)</u> opening of a file in the first computer entity <u>b)</u>

<u>opening a debugging port in the first computing entity</u>. (see Yan paragraph [0054], lines

1-14: metrics that reflect configuration state; a metric may change with time, this

requiring a new value to be stored; specification paragraph [0027] discloses opening

and reading a file used to modify security environment)

Regarding Claim 34, Yan discloses the method of claim 31, wherein the behavior

parameter comprises opening a debugging port in the first computer entity. (see Yan

paragraph [0054], lines 1-14: metrics that reflect configuration state; a metric may

change with time, this requiring a new value to be stored; specification paragraph [0027]

discloses debugging port used to modify security environment)

Regarding Claim 35, Yan discloses the method of claim 31, wherein the trust

message. (see Yan paragraph [0017], lines 1-6: establishing and maintaining trust

between a trustee and a trustor; generating metrics of a trustor and comparing to

current metrics of a trustee)   Yan does not specifically disclose a period of time

whereby a secret is valid.   However, Qui further discloses data to inform the first

computer entity of the first period of time over which the first secret is valid. (see Qui

paragraph [0040], lines 1-7; paragraph [0021], lines 8-11: expiration timer for certificate

information)

It would have been obvious to one of ordinary skill in the art to modify Yan for a

period of time whereby a secret is valid as taught by Qui. One of ordinary skill in the art

would have been motivated to employ the teachings of Qui for the generation,

transmission, and updating of certificate information when the number of devices is

large. (see Qui paragraph [0007], lines 7-12)

Regarding Claim 36, Yan discloses the method of claim 31, further comprising:

retransmitting the trust message from the second computer entity to the first computer

entity, the retransmitted trust message including a) a second secret that is different than

the first secret,

Yan does not specifically disclose a period time whereby a secret is valid.

However, Qui further discloses data to inform the first computer entity of a second

period of time over which the second secret is valid. (see Qui paragraph [0040], lines 1-

7; paragraph [0021], lines 8-11: expiration timer for certificate information)

It would have been obvious to one of ordinary skill in the art to modify Yan for a

trust message further including an expiration time as taught by Qui. One of ordinary

skill in the art would have been motivated to employ the teachings of Qui for the

generation, transmission, and updating of certificate information when the number of

devices is large. (see Qui paragraph [0007], lines 7-12)

Regarding Claim 37, Yan discloses the method of claim 30, wherein the can-attest

message is transmitted from the first computer entity to the server in an unencrypted

format. (see Yan paragraph [0064], lines 1-17: initial attestation transmitted via

message flow between trustor and trustee to establish trust relationship)

Regarding Claim 38, Yan discloses the method of claim 30, wherein each of the can-attest message and the attestation-wanted message is configured to preclude a) containing or b) using a cryptographic key. (see Yan paragraph [0064], lines 1-17: initial attestation transmitted via message flow between trustor and trustee to establish trust relationship)

6.   Claims **16, 28** are rejected under 35 U.S.C. 103(a) as being unpatentable over **Yan** in view of **Grawrock** (US PGPUB No. **20040117625**).

Regarding Claims 16, Yan discloses the method of claim 31 wherein the second computer entity creates the trust message by application of that automatically produces the trust message in an appropriate format that is accessible to the first computer entity. (see Yan paragraph [0059], lines 6-12: generate formatted attestation information) Yan does not specifically disclose a sealing function.

However, Grawrock discloses wherein a sealing function. (see Grawrock paragraph [0018], lines 12-16; paragraph [0025], lines 1-7; paragraph [0026], lines 1-6: seal/unseal trusted operation utilized)

It would have been obvious to one of ordinary skill in the art to modify Yan to perform a seal operation as taught by Grawrock. One of ordinary skill in the art would have been motivated to employ the teachings of Grawrock to provide local users and

remote computing devices an efficient and easier method for the completion of trusted

operations. (see Grawrock paragraph [0002], lines 7-14)

Regarding Claims 28, Yan discloses the method of claim 19 wherein the first computer

entity consumes the trust message by application of same that automatically extracts

the shared secret and other relevant information from such trust attestation message for

use by the first computer entity. (see Yan paragraph [0060], lines 6-9: extracts

attestation information for processing)

Yan does not specifically disclose whereby an unsealing function.

However, Grawrock discloses wherein an unsealing function. (see Grawrock paragraph

[0018], lines 12-16; paragraph [0025], lines 1-7; paragraph [0026], lines 1-6: seal/unseal

trusted operation utilized)

It would have been obvious to one of ordinary skill in the art to modify Yan perform

an unseal operation within a trusted computing environment as taught by Grawrock.

One of ordinary skill in the art would have been motivated to employ the teachings of

Grawrock to provide local users and remote computing devices an efficient and easier

method for the completion of trusted operations. (see Grawrock paragraph [0002], lines

7-14)

8.   Claim **30** is rejected under 35 U.S.C. 103(a) as being unpatentable over **Yan** in

view of **Pinkas**.

Regarding Claims 30, Yan discloses a method of establishing trust between a first

computer entity and <u>server</u>, the method comprising: first computer entity constructing

the attestation message, the first computer entity <u>seeking a granting of trust from the</u>

<u>server by</u> sending a can-attest message to the <u>server</u>, the can-attest message stating

that the first computer entity can send an attestation message but that the first computer

entity would like to know from the <u>server</u> whether such an attestation message is

required, and if so any requirements that such <u>server</u> has with regard to such attestation

message, whereby the <u>server</u> sends an attestation-wanted message to the first

computer entity in response to the can-attest message, the attestation-wanted message

stating that the <u>server</u> does in fact require an attestation message from the first

computer entity and that the attestation message as sent by the first computer entity

must adhere to certain requirements as defined in such attestation-wanted message,

the first computer entity thereafter sending the attestation message in accordance with

the requirements stated in the attestation-wanted message. (see Yan paragraph

[0064], lines 1-17: initial attestation transmitted via message flow between trustor and

trustee to establish trust relationship; Yan discloses a two step process for attestation

messages (message flows 302, 304 and message flows 306, 308))

Yan does not specifically disclose a first entity sends an initial attestation.

However, Pinkas discloses or can-attest message or an initial message transmitted

from a first entity. (see Pinkas col 4, line 54 - col. 5, line 2: requestor (trustee, first

entity) addresses an attestation request or query to the authority; server transmits

attestation to requestor subject)

It would have been obvious to one of ordinary skill in the art to modify Yan-Qui whereby a first entity sends an initial attestation message as taught by Pinkas.   One of ordinary skill in the art would have been motivated to employ the teachings of Pinkas to provide flexible and simplicity in the establishment of attestation while provide security with the usage of a check datum. (see Pinkas col. 4, lines 44-53)

Conclusion

Any inquiry concerning this communication or earlier communications from the examiner should be directed to CARLTON V. JOHNSON whose telephone number is (571)270-1032. The examiner can normally be reached on Monday thru Friday , 8:00 - 5:00PM EST.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Nasser Moazzami can be reached on 571-272-4195. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

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USPTO Customer Service Representative or access to the automated information

system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.

/Eleni A Shiferaw/   Carlton V. Johnson
Primary Examiner, Art Unit 2436   Examiner
                                                        Art Unit 2436

CVJ
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